

Sirindhorn International Institute of Technology

Thammasat University at Rangsit

School of Information, Computer and Communication Technology

## TCS455: Midterm Examination (Set I)

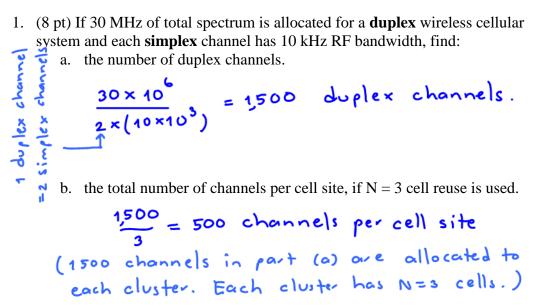
Solution

COURSE: TCS455 (Mobile Communications)DATE: December 22, 2009SEMESTER: 2/2009INSTRUCTOR:Dr. Prapun SuksompongTIME: 9:00-12:00PLACE: BKD 26XX

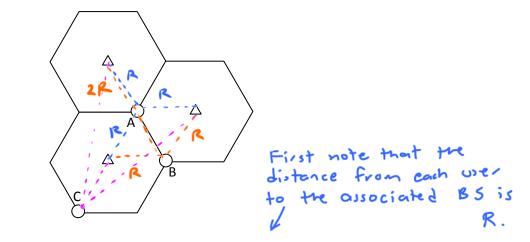
Name Prapun	ID	555
	Seat	

## **Instructions:**

- 1. Including this cover page, this exam has 9 pages.
- 2. Read these instructions and the questions carefully.
- 3. Closed book. Closed notes.
- 4. Basic calculators are permitted, but borrowing is not allowed.
- 5. Allocate your time wisely. Some easy questions give many points.
- 6. Do not cheat. The use of communication devices including mobile phones is prohibited in the examination room.
- 7. (1 pt) Write your **first name and the <u>last three digits</u> of your ID** on each page of your examination paper, starting from page 2.
- 8. Your score depends strongly on your explanation of your answer. If the explanation is incomplete, zero score may be given even when the final answer is correct.
- 9. Do not panic.
- 10. Dr. Prapun will visit each exam room regularly. In general, there is no need to ask the proctor to call for Dr. Prapun.



2. (15 pt) In this question, we consider the SIR value when the cluster size N = 1. Sectoring is not used. Suppose this cellular system has only three base stations. They are marked by triangles located at the centers of three cells in the figure below. Assume that they transmit the same power level.

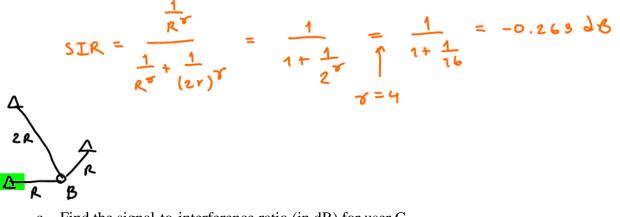


User A (mobile station A), user B (mobile station B), and user C (mobile station C) are currently associated with the lower-left base station. The locations of these users are marked by the circles in the figure above. Assume a path loss exponent of  $\gamma = 4$ . Do not approximate distance values.

a. Find the signal-to-interference ratio (in dB) for user A.

The distance from A to all BS = R.
So, SIR = $\frac{P_r}{ZP_i} = \frac{1/R^2}{R^2 + \frac{1}{R^2}} = \frac{1}{2} = -3 dB$ $A_R$
$A_{0} = \frac{2P_{i}}{R^{2}} \uparrow \frac{1}{R^{2}} \downarrow \frac{1}{R^{2}}$
ARA Note that we have only two intererferers.
R intererferers.
Page 2 of 9
Page 2 of 9 Note also that the signals from Bond Cwill
not interfere with the signal from A becare they use different fregs.

b. Find the signal-to-interference ratio (in dB) for user B.



c. Find the signal-to-interference ratio (in dB) for user C.

$$2^{2} + \sqrt{3}^{2} R = \sqrt{7} R$$

$$2^{2} + \sqrt{3}^{2} R = \sqrt{7} R$$

$$SIR = \frac{1}{R^{2}} = \frac{\sqrt{7}}{2} = \frac{19}{2} = 13.9 d6$$

$$2 \times \frac{1}{2} R = \frac{\sqrt{7}}{2} = \frac{1}{2} = 13.9 d6$$

$$(\sqrt{7}R)^{7} r = 4$$

3. (15 pt) A cellular service provider decides that its users can tolerate a signal-to-interference ratio of 20 dB in the worst case. Assume a path loss exponent of γ = 4. You may approximate the distances between the users and the co-channel base stations by the center-to-center co-channel distance D.
 a. Find the optimal value of cluster size N for omnidirectional antennas.

since 
$$\frac{r}{k} = \frac{r}{k} = \frac{r}{k} = \frac{1}{k} \frac{p}{k} = \frac{1}{k} \left(\frac{p}{k}\right)^{2} = \frac{1}{k} \left(\sqrt{3N}\right)^{2} = \frac{1}{2} q N^{2} \geq 10$$
  
 $i = 1$   $-k = 6$  for omnidirectional  $r = 4$ 

۵

$$\Rightarrow N = \sqrt{\frac{100 \text{ K}}{9}} = 10 \times \sqrt{\frac{2}{3}} = 8.165 \longrightarrow \text{Optimal N = 9}$$

b. Find the optimal value of cluster size N for  $120^{\circ}$  sectoring.

$$N \ge \sqrt{\frac{100 \text{ K}}{9}} = 4.714 \longrightarrow \text{Optimal N} = 7$$
  
 $K = 2$   
for 120°  
sectoring

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c. Find the optimal value of cluster size N for 60° sectoring.  

$$N \gtrsim \sqrt{\frac{100 \text{ K}}{9}} = \frac{10}{3} = 3.33 \longrightarrow \text{Optimal N} = 4$$
  
 $\kappa = 1 \text{ for}$   
 $60^{\circ}$   
sectoring  
 $m = 2$ 

4. (5 pt) There are 1000 users subscribed to a cellular system. The call request rate for each user is 2 call requests per week. For each call, the average call duration is 1 min. If the system has only two channels and it is used to support the whole 1000 users, what is the blocking probability?

$$\lambda = 2 \times 1000 \quad \frac{requests}{week} \times \frac{1}{p} \frac{day}{days} \times \frac{1}{24} \frac{hr}{hrs.} \qquad \frac{1}{60} \frac{hr}{min}$$

$$\frac{1}{m} = 1 \quad \min n$$

$$\frac{1}{m} = 2 \quad 0.19 \quad \text{s} \quad \frac{1}{p} \quad \frac{1}{p} \quad \frac{4^{2}}{2} = 0.016 = 1.6\%$$

$$m = 2$$

$$5. \text{ A function ErlangB}(m, A) \text{ is defined by} \qquad \text{If } you \quad \text{don't remember}$$

$$\text{tre } formula, \quad \text{it is}$$

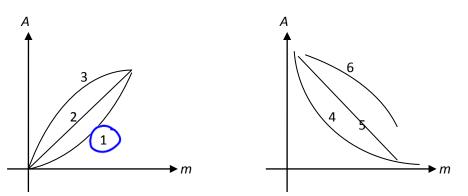
$$\text{ErlangB}(m, A) = \frac{\frac{A^{m}}{m!}}{\sum_{k=0}^{m} \frac{A^{k}}{k!}} \qquad \text{given in the next}$$

a. (4 pt) Why is this function useful? How can we use it in designing cellular system? Your score depends on the completeness, correctness, and clarity of your answer.

For a cell (or a sector) in a cellular system which has m channels and the amount of traffic is A Erlang, this function is directly used to determine the probability Pb that call requests will be blocked by the system because all channels are currently used. The amount of traffic (A) can be found by the product of the call request rate and the average call duration. When we design a cellular system, the blocking probability Pb should be less than some pre-determined value. In which case, the funtion above can be used to suggest the minimum number of channels per cell (or sector). If we already know the number of channels per cell (or sector) of the system, the function can also be used to determined how many users the system can support as well.

Blocking probability Pb and the SIR values determine the quality of service for cellular users. We can use Pb to trade with SIR and vice versa.

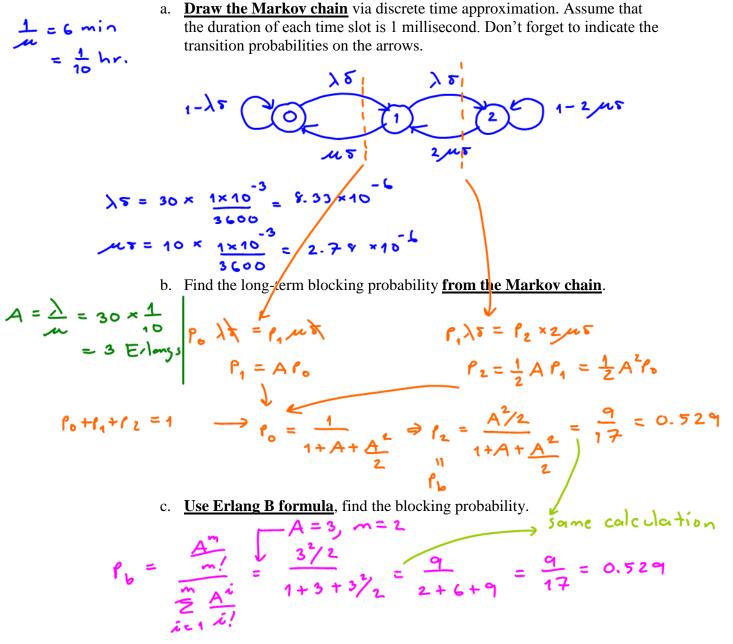
ID A А m b. (2 pt) Compare ErlangB(1000, 2000) and ErlangB(1000, 2001). Which one is larger? more Erlang ⇒ mure traffic => larger Ph c. (2 pt) Compare ErlangE((1000, 2000) and ErlangB(1001, 2000). Which one is larger? less Xichannels => larger Pb d. (1 pt) Suppose  $ErlangB(1000, A_1) = ErlangB(2000, A_2) = ErlangB(3000, A_3) = 0.05$ Let  $A_4 = A_1 + A_2$ . Compare  $A_3$  and  $A_4$ . Which one is larger? 3,000 channels are divided into two groups of size 1,000 and 3000. So, the trunking efficiency should be higher for the 3 channels care; that is it should be oble to support traffic for fue same blocking probability. e. (1 pt) Consider the relation ErlangB(m, A) = 0.005. Which plot below best represents the relationship between *m* and *A*?



Hint: No calculation is needed. Think about what the formula says.

we know that as we increase my the increased & channel allows us to support more traffic (A) - So, the graph should be on increasing function. Now, suppose we increase m by 2x, we know, by the increased trunking efficiency, that A should be increased by more Page 5 of 9 than 2x.

6. (15 pt) Consider a system which has 2 channels. We would like to find the blocking probability via the Markov chain method. Assume that the total call request rate is 30 calls per hour and the average call duration is 6 mins.



7. (5 pt) What do we mean when we say a Markov chain is memoryless? Your score depends on the completeness, correctness, and clarity of your answer.

The future is conditionally independent of the past history, given the present state. In other words, the entire past history is summarized in the present state. We need no memory of the past history (only the present) to probabilistically predict the future.

This memoryless property is also referred to as the Markov property, after A.A. Markov who first described and studied this class of processes in the early 1900s.

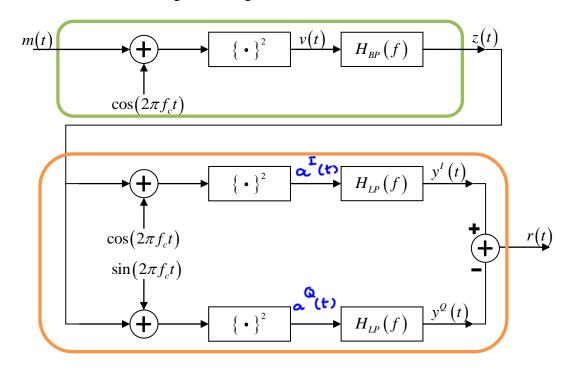
8. (14 pt) Complete the following M/M/m/m queue description with the following terms:

(I) Bernoulli (IV) Gaussian	(II) binomial (V) geometric	(III) exponential (VI) Poisson	We haven't a these interva	assumed that
assumptions <b>foisson(A)</b> <b>expone(B)</b> consider nor intervals are adjacent call In order to a slotted time our focus fro In the limit,	random vari n-overlapping time int i.i.d. <u>Po(©) on</u> requests are i.i.d nalyze or simulate the where the length of e. om continuous-time M only one of the two e	est process is modeled (2) the call durations ables. For the call re- tervals, the numbers random variab $cx \in C$ and $c$ radius e system described all ach time slot is small Markov chain to discri- vents can happen dur	ons. Two important d by a/an s are assumed to be i.i.d quest process, if we of call requests in these les. The times between andom variables.	d. e i 1.
<u>Ber</u> (D) U number of ca	of new call requests $\frac{1}{2}$ random variall requests during <i>n</i> sable because it is a su	iables. <u>In which case</u> slots, we will get a/ar im of i.i.d. <u>Ber</u>	, if we count the total $h = bi(E) \circ m i \circ a$	noulli assump.
number of sl limit, the num by a <u>bia</u> foisso(C)	<b>b)<u>mia</u></b> random random vari	ill increase as the slot in the time interval <i>I</i> variable before will able.	ts get smaller. In the which we approximate approach a/an f any thing is	a discrete
these number random varia	ables can be thought of the num	of as discrete counter	rparts of the i.i.d.	,

Put your answers in the box below. Some terms above are used more than once. Some terms are not used.

(A)	VI
(B)	Ħ
(C)	<b>↓</b> Ľ
(D)	Н
(E)	Ħ
(F)	<
(G)	Ħ

9. Suppose  $m(t) \xrightarrow{\mathcal{F}} M(f)$  is bandlimited to *W*, i.e., |M(f)| = 0 for |f| > W. Consider the following block diagram:



where  $f_c \gg W$ ,

Let  $w_c = 2\pi f_c$  $H_{BP}(f) = \begin{cases} 1, & |f - f_c| \le W \\ 1, & |f + f_c| \le W \\ 0, & \text{otherwise} \end{cases}$  and  $H_{LP}(f) = \begin{cases} 1, & |f| \le W \\ 0, & \text{otherwise.} \end{cases}$ 

a. (5 pt) Give a simplified formula for 
$$z(t)$$
.  

$$w(t) = (m(t) + \cos(w_{c}t))^{2} = m^{2}(t) + 2m(t)\cos(w_{c}t) + \cos^{2}(w_{c}t)$$

$$BPF$$

$$w(t) = \frac{1}{2}\cos(2w_{c}t)$$

$$BPF = PFF$$

 $z(t) = 2m(t)cos(w_ct)$ 

b. (5 pt) Give a simplified formula for 
$$r(t)$$
.  

$$\alpha^{T}(t) = (Z(t) + \cos \omega_{c} t)^{2} = (2m(t) \cos \omega_{c} t + \cos \omega_{c} t)^{2} = \cos^{2}\omega_{c} t (2m(t)+t)^{e}$$

$$= (\frac{1}{2} + \frac{1}{2}\cos(2\omega_{c} t)) (4m^{2}(t) + 4m(t) + 1)$$

$$= (\frac{1}{2} + \frac{1}{2}\cos(2\omega_{c} t)) (4m^{2}(t) + 4m(t) + 1)$$

$$= (2(t) + \sin \omega_{c} t)^{2} = (2m(t) \cos \omega_{c} t + \sin \omega_{c} t)^{2}$$

$$= 4m^{2}(t) \cos^{2}\omega_{c} t + 4m(t) \cos \omega_{c} t + \sin \omega_{c} t)^{2}$$

$$= 4m^{2}(t) \cos^{2}\omega_{c} t + 4m(t) \cos \omega_{c} t + \sin \omega_{c} t)$$

$$= 2m^{2}(t) + 2m^{2}(t) \cos(2\omega_{c} t) + m(t) \sin(2\omega_{c} t) + \frac{1}{2} - \frac{1}{2}\cos(2\omega_{c} t)$$

$$LIFF$$

$$= LIFF \{2m^{2}(t)\} + \frac{1}{2}$$

$$r(t) = \sqrt{T}(t) - \sqrt{G}(t) = 2m(t)$$

$$c. (2pt) Suppose  $M(f) = \begin{cases} 1, |f| \le W \\ 0, \text{ otherwise.} \end{cases}$$$

Plot |V(f)|. Your score depends strongly on your explanation and how accurate your plot is.

$$From (a), \forall r(t) = m^{t}(t) + 2 m(t) \cos \left(m_{c}t\right) + \frac{1}{2} + \frac{1}{2} \cos \left(2m_{c}t\right)$$

$$\int_{1}^{2} d\tau = 2m$$

$$\int_{1}^{2} d\tau = 2m$$